

HP-SEE Valgrind Usage

www.hp-see.eu

HP-SEE

High-Performance Computing Infrastructure for South East Europe's Research Communities

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Introduction



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- Valgrind is a suite of command line tools both for debugging and profiling codes on Linux system.
- Includes a set of production-quality tools
 - Memcheck memory error detector
 - Cachegrind cache and branch-prediction profiler
 - Callgrind call-graph generating extension to Cachegrind
 - Massif heap profiler
 - Helgrind thread error detector
- Ease of use:
 - Valgrind uses dynamic binary instrumentation no need to modify, recompile or relink your applications.
 - Simply prefix your command line with valgrind and everything works.

Why you should use it



- Valgrind helps solving issues with dynamic memory allocation and errors associated with it:
 - Automatically detect many memory management and threading bugs, saving hours of debugging time.
 - Valgrind tools allow very detailed profiling to help find bottlenecks in your programs, often resulting in program speed-up.
 - Valgrind works with programs written in any language.
 - □ Valgrind works with MPI: Open-MPI and MVAPICH/MVAPICH2
- Downsides:
 - Large overhead
 - Programs run significantly more slowly under Valgrind. Depending on which tool you use, the slowdown factor can range from 5 – 100.
 - Measurements may not be absolutely accurate

Common Errors



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- Use of uninitialized memory
- Reading/writing memory after it has been freed
- Reading/writing off the end of allocated blocks
- Reading/writing inappropriate areas on the stack
- Memory leaks where pointers to allocated blocks become lost
- Mismatched use of malloc/new/new[] vs free/delete/ delete[]

Compiling



- When testing for memory problems It is recommended to compile the code with both the debugging options -O0 (no optimization) and -g (debugging information).
 - □ \$(CC) filecode.c –g –O0 –o fileprog.x
 - □ \$(FC) filecode.f90 -g -O0 -o fileprog.x
 - \$(CXX) filecode.cpp -g -O0 -fno-inline -o fileprog.x
- The –fno-inline flag avoids the inlining of functions into the main program and makes it easier to see the function-call chain.
- Using Valgrind with code that has been compiled with optimisation options could give incorrect results.
- These examples can also be applied using the MPI compiler wrappers.

Memcheck: Memory Error Checker



- □ Aimed primarily at Fortran, C and C++ programs.
- All reads and writes of memory are checked, and calls to malloc/new/free/delete are intercepted. Will report if:
 - Accesses memory it shouldn't (not yet allocated, freed, past the end of heap blocks, inaccessible areas of the stack).
 - Uses uninitialized values in dangerous ways.
 - Leaks memory.
 - Does bad frees of heap blocks (double frees, mismatched frees).
 - Passes overlapping source and destination memory blocks to memcpy() and related functions.
- Memcheck reports these errors as they occur, giving the source line number, and also a stack trace of the functions called to reach that line.
- □ Memcheck runs programs $10-30 \times$ slower than normal.

Memcheck usage Uninitialized Memory



```
[josipjakic@ui Valgrind]$ valgrind --tool=memcheck ./uninit_memory
1 #include <stdlib.h>
                              ==32126== Memcheck, a memory error detector
2 int main() {
                              ==32126== Copyright (C) 2002-2009, and GNU GPL'd, by Julian Seward et al.
3
                              ==32126== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info
4
   int p, t;
                               ==32126== Command: ./uninit_memory
5
                              ==32126==
6
   if (p == 5) /* Error */
                               ==32126== Conditional jump or move depends on uninitialised value(s)
     t = p + 1;
7
                              ==32126== at 0x400450: main (uninit_memory.c:6)
                              ==32126==
8
   return 0;
9 }
                              ==32126==
                              ==32126== HEAP SUMMARY:
                                             in use at exit: 0 bytes in 0 blocks
                               ==32126==
                              ==32126== total heap usage: 0 allocs, 0 frees, 0 bytes allocated
p is uninitialized and
                              ==32126==
                               ==32126== All heap blocks were freed -- no leaks are possible
may contain garbage,
                              ==32126==
resulting in an error if
                              ==32126== For counts of detected and suppressed errors, rerun with: -v
                              ==32126== Use --track-origins=yes to see where uninitialised values come from
used to determine
                              ==32126== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 4 from 4)
branch-outcome or
                              [josipjakic@ui Valgrind]$
memory address
```

(ex: a[p] = y)

Memcheck usage Invalid Read/Write



<pre>1 #include <stdlib.h> 2 int main() { 3 4 int *p, i, a; 5 6 m mpllag(10*sizesf(int)); </stdlib.h></pre>	[josipjakic@ui Valgrind]\$ valgrindtool=memcheck ./invalid_read_write ==13028== Memcheck, a memory error detector ==13028== Copyright (C) 2002-2009, and GNU GPL'd, by Julian Seward et al. ==13028== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info ==13028== Command: ./invalid_read_write ==13028==	High-Perform for South East
<pre>6 p = malloc(10*sizeof(int)); 7 p[11] = 1; /* write */ 8 a = p[11]; /* read */ 9 free(p); 10 return 0; 11 }</pre>	<pre>==13028== Invalid write of size 4 ==13028== at 0x4004F6: main (invalid_read_write.c:7) ==13028== Address 0x4c3b06c is 4 bytes after a block of size 40 alloc'd ==13028== at 0x4A05E1C: malloc (vg_replace_malloc.c:195) ==13028== by 0x4004E9: main (invalid_read_write.c:6) ==13028== ==13028== Invalid read of size 4 ==13028== at 0x400504: main (invalid_read_write.c:8)</pre>	
Attempting to read/ write from address (p+sizeof(int)*11) which has not been allocated.	<pre>==13028== Address 0x4c3b06c is 4 bytes after a block of size 40 alloc'd ==13028== at 0x4A05E1C: malloc (vg_replace_malloc.c:195) ==13028== by 0x4004E9: main (invalid_read_write.c:6) ==13028== ==13028== in use at exit: 0 bytes in 0 blocks ==13028== in use at exit: 0 bytes in 0 blocks ==13028== total heap usage: 1 allocs, 1 frees, 40 bytes allocated ==13028== ==13028== All heap blocks were freed no leaks are possible ==13028== ==13028== = ==13028== For counts of detected and suppressed errors, rerun with: -v ==13028== ERROR SUMMARY: 2 errors from 2 contexts (suppressed: 4 from 4 [josipjakic@ui Valgrind]\$</pre>	·)

Memcheck usage Invalid Free



<pre>1 #include <stdlib.h> 2 3 int main() { 4 5 int *p, i; 6 p = malloc(10*sizeof (int)); 7 for(i = 0;i < 10;i++) 8 p[i] = i; 9 free(p); 10 free(p); /* Error */ 11 return 0; 12 }</stdlib.h></pre>	<pre>[josipjakic@ui Valgrind]\$ valgrindtool=memcheck ./invalid_free ==24100== Memcheck, a memory error detector ==24100== Copyright (C) 2002-2009, and GNU GPL'd, by Julian Seward et al. ==24100== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info ==24100== Command: ./invalid_free ==24100== ==24100== Invalid free() / delete / delete[] ==24100== at 0x4A05A31: free (vg_replace_malloc.c:325) ==24100== by 0x400527: main (invalid_free.c:9) ==24100== Address 0x4c3b040 is 0 bytes inside a block of size 40 free'd ==24100== at 0x4A05A31: free (vg_replace_malloc.c:325) ==24100== by 0x40051E: main (invalid_free.c:8) ==24100== ==24100== ==24100==</pre>
Valgrind checks the address passed to the free() call and sees that it has already been freed.	<pre>==24100== HEAP SUMMARY: ==24100== in use at exit: 0 bytes in 0 blocks ==24100== total heap usage: 1 allocs, 2 frees, 40 bytes allocated ==24100== ==24100== All heap blocks were freed no leaks are possible ==24100== ==24100== For counts of detected and suppressed errors, rerun with: -v ==24100== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 4 from 4) [josipjakic@ui Valgrind]\$</pre>

Memcheck usage Invalid Call Parameter



<pre>1 #include <stdlib.h> 2 #include <unistd.h> 3 4 int main() { 5 int *p; 6 7 p = malloc(10); 8 read(0, p, 100); /* err */ 9 free(p); 10 return 0; 11 }</unistd.h></stdlib.h></pre>	<pre>[josipjakic@ui Valgrind]\$ valgrindtool=memcheck ./invalid_call_param ==18300== Memcheck, a memory error detector ==18300== Copyright (C) 2002-2009, and GNU GPL'd, by Julian Seward et al. ==18300== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info ==18300== Command: ./invalid_call_param ==18300== ==18300== at 0x3351AC52A0:read_nocancel (in /lib64/libc-2.5.so) ==18300== by 0x400550: main (invalid_call_param.c:7) ==18300== Address 0x4c3b04a is 0 bytes after a block of size 10 alloc'd ==18300== by 0x400539: main (invalid_call_param.c:6) ==18300== 12345678901234567890</pre>
read() tries to read 100 bytes from stdin and place the results in p but the bytes after the firs 10 are unaddressable.	<pre>==18300== ==18300== HEAP SUMMARY: ==18300== in use at exit: 0 bytes in 0 blocks ==18300== total heap usage: 1 allocs, 1 frees, 10 bytes allocated ==18300== ==18300== All heap blocks were freed no leaks are possible ==18300== ==18300== For counts of detected and suppressed errors, rerun with: -v ==18300== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 4 from 4) [josipjakic@ui Valgrind]\$</pre>

Memcheck usage Leak Detection



<pre>1 #include <stdlib.h> 2 3 int main() { 4 int *p, i; 5 p = malloc(5*sizeof(int)); 6 for(i = 0;i < 5;i++) 7 p[i] = i; 8 return 0; 9 }</stdlib.h></pre>	<pre>[josipjakic@ui Valgrind]\$ valgrindleak-check=yestool=memcheck ./memory_leak ==325== Memcheck, a memory error detector ==325== Copyright (C) 2002-2009, and GNU GPL'd, by Julian Seward et al. ==325== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info ==325== Command: ./memory_leak ==325== ==325== ==325== ==325== in use at exit: 20 bytes in 1 blocks ==325== total heap usage: 1 allocs, 0 frees, 20 bytes allocated ==325==</pre>
20 unfreed blocks at routine exit – memory leak.	<pre>==325== 20 bytes in 1 blocks are definitely lost in loss record 1 of 1 ==325== at 0x4A05E1C: malloc (vg_replace_malloc.c:195) ==325== by 0x4004A9: main (memory_leak.c:5) ==325== ==325== LEAK SUMMARY: ==325== definitely lost: 20 bytes in 1 blocks ==325== indirectly lost: 0 bytes in 0 blocks ==325== possibly lost: 0 bytes in 0 blocks ==325== still reachable: 0 bytes in 0 blocks ==325== suppressed: 0 bytes in 0 blocks ==325== = ==325== For counts of detected and suppressed errors, rerun with: -v ==325== ERROR SUMMARY: 1 errors from 1 contexts (suppressed: 4 from 4) [josipjakic@ui Valgrind]\$</pre>

Cachegrind: Cache profiler



- Performs detailed simulation of I1, D1 and L2 caches
- Can accurately pinpoint the sources of cache misses in your code. It identifies for each line of source code the number of:
 - Cache misses
 - Memory references
 - Instructions executed
- Provides per-function, per-module and whole-program summaries.
- Useful for programs written in any language.
- Performance hit is about a 20-100x slowdown.

Cachegrind usage (1/2)

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#include <stdio.h> 21 #define N 1000 double array sum(double a[][N]) { 2 22 з 23 4 double array sum(double a[][N]); 24 int i,j; 5 double s; 25 int main(int argc, char **argv) { 6 26 7 27 s=0: 8 double a[N][N]; 28 for (i=0;i<N;i++) Read 2D 9 int i,j; 29 for (j=0;j<N;j++) 30 10 s += a[i][j]; Array for (i=0;i<N;i++) { 31 11 Fill 2D for (j=0;j<N;j++) { 32 12 return s; 13 a[i][j] = 0.01;33 - 1 Array 14 1 15 ъ 16 17 printf("sum = %10.3f n", array sum(a));18 19 return 0; 20

Array size is 1,000 x 1000 x 8 bytes = 8Mb
 32kB L1i and 32kB L1d
 4096kB L2

Cachegrind usage (2/2)

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```
[josipjakic@ui Valgrind]$ gcc -O2 -g -o loops-fast loops-fast.c
[josipjakic@ui Valgrind]$ valgrind --tool=cachegrind ./loops-fast
==23430== Cachegrind, a cache and branch-prediction profiler
==23430== Copyright (C) 2002-2009, and GNU GPL'd, by Nicholas Nethercote et al.
==23430== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info
==23430== Command: ./loops-fast
==23430==
sum = 10000.000
==23430==
==23430== I refs:
                     10,122,886
==23430== I1 misses:
                           847
==23430== L2i misses:
                           846
==23430== I1 miss rate:
                           0.00%
                           0.00%
==23430== L2i miss rate:
==23430==
==23430== D refs:
                      2,041,972 (1,029,938 rd + 1,012,034 wr)
==23430== D1 misses:
                         251,113 (125,846 rd + 125,267 wr)
==23430== L2d misses:
                         251,047 (125,785 rd + 125,262 wr)
                           12.2% (
                                     12.2%
                                                   12.3%)
==23430== D1 miss rate:
                                              +
==23430== L2d miss rate:
                           12.2% (
                                     12.2%
                                                   12.3%)
                                              +
==23430==
==23430== L2 refs:
                       251,960 (126,693 rd + 125,267 wr)
==23430== L2 misses:
                         251,893 (126,631 rd + 125,262 wr)
                            2.0% (
                                     1.1% +
                                                 12.3%)
==23430== L2 miss rate:
[josipjakic@ui Valgrind]$
```

Callgrind: Callgraphs+Cachegrind Info



- Is an extension that provides all the info Cachegrind yields
- Provides callgraph information.
- Kcachegrind is a separately available tool for visualisation for both Callgrind and Cachegrind output data

Callgrind usage (1/3)

```
[josipjakic@ui Valgrind]$ valgrind --tool=callgrind --simulate-cache=yes ./loops-fast
==19141== Callgrind, a call-graph generating cache profiler
==19141== Copyright (C) 2002-2009, and GNU GPL'd, by Josef Weidendorfer et al.
==19141== Using Valgrind-3.5.0 and LibVEX; rerun with -h for copyright info
==19141== Command: ./loops-fast
==19141==
==19141== For interactive control, run 'callgrind_control -h'.
sum = 10000.000
==19141==
==19141== Events : Ir Dr Dw I1mr D1mr D1mw I2mr D2mr D2mw
==19141== Collected : 10122883 1029478 1012494 847 125838 125275 846 125777 125270
==19141==
==19141== I refs:
                     10,122,883
==19141== I1 misses:
                           847
==19141== L2i misses:
                           846
==19141== I1 miss rate:
                           0.0%
==19141== L2i miss rate:
                           0.0%
==19141==
                      2,041,972 (1,029,478 rd + 1,012,494 wr)
==19141== D refs:
==19141== D1 misses:
                         251,113 (125,838 rd + 125,275 wr)
==19141== L2d misses:
                         251,047 (125,777 rd + 125,270 wr)
==19141== D1 miss rate:
                           12.2\%(12.2\%+12.3\%)
                           12.2% ( 12.2% +
                                                 12.3%)
==19141== L2d miss rate:
==19141==
==19141== L2 refs:
                       251,960 (126,685 rd + 125,275 wr)
==19141== L2 misses:
                         251,893 (126,623 rd + 125,270 wr)
==19141== L2 miss rate:
                           2.0% ( 1.1% +
                                                12.3%)
[josipjakic@ui Valgrind]$
```

Tuning and Optimization of HPC Application – Institute of Physics Belgrade, Friday 01 June 2012

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Callgrind usage (2/3)



Cachegrind saves output to a file *callgrind.out.<pid>* Use callgrind_annotate to parse this file for detailed information

[josipjakic@ui Valgrind]\$ callgrind_annotate callgrind.out.19141 _____ Profile data file 'callgrind.out.19141' (creator: callgrind-3.5.0) I1 cache: 32768 B, 64 B, 8-way associative D1 cache: 32768 B, 64 B, 8-way associative L2 cache: 4194304 B, 64 B, 16-way associative Timerange: Basic block 0 - 2024406 **Trigger: Program termination** Profiled target: ./loops-fast (PID 19141, part 1) Events recorded: Ir Dr Dw I1mr D1mr D1mw I2mr D2mr D2mw Ir Dr Dw I1mr D1mr D1mw I2mr D2mr D2mw Events shown: Event sort order: Ir Dr Dw I1mr D1mr D1mw I2mr D2mr D2mw Thresholds: 9900000000 Include dirs: User annotated: Auto-annotation: off

Callgrind usage (3/3)

Dw I1mr D1mr D1mw I2mr D2mr D2mw Ir Dr 10,122,886 1,029,478 1,012,494 847 125,838 125,275 846 125,777 125,270 PROGRAM TOTALS Dw I1mr D1mr D1mw I2mr D2mr D2mw file:function Ir Dr 6,007,017 4 1,000,004 2 2 125,001 2 2 125,001 loops-fast.c:main [/home/josipjakic/Valgrind/loops-fast] 4,005,003 1,000,001 0 0 125,001 0 0 125,001 . loops-fast.c:array_sum [/home/josipjakic/Valgrind/loopsfast] 23,333 7,843 3,388 12 163 5 12 157 4 ???:do lookup x [/lib64/ld-2.5.so] [josipjakic@ui Valgrind]\$

Callgrind can be used to find performance problems that are not related to CPU cache

□ What lines eat up most instructions (CPU cycles, time)

What system/math/lib functions are called and what is their cost

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Massif: Heap Profiler

- Performs detailed profiling by taking regular snapshots of a program's heap.
- Produces a graph showing heap usage over time
- Massif runs programs about 20× slower than normal.

Helgrind: Thread Debugger

- Finds data races in multithreaded programs.
- Looks for memory locations which are accessed by more than one [POSIX p-]thread
- It is useful for any program that uses pthreads.
- Experimental tool





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Valgrind is freely available from:

http://www.valgrind.org